

IMFM

INSTITUTE OF MATHEMATICS, PHYSICS AND MECHANICS
JADRANSKA 19, 1000 LJUBLJANA, SLOVENIA

Preprint series

Vol. 48 (2010), 1114

ISSN 2232-2094

ON THE 2-RESONANCE OF
FULLERENES

Tomáš Kaiser Matěj Stehlík

Riste Škrekovski

Ljubljana, February 26, 2010

On the 2-resonance of fullerenes*

Tomáš Kaiser[†] Matěj Stehlík[‡] Riste Škrekovski[§]

February 24, 2010

Abstract

We show that every pair of hexagons in a fullerene graph satisfying the isolated pentagon rule (IPR) forms a resonant pattern. This solves a problem raised by Ye *et al.* [SIAM J. Discrete Math. 23(2):2009, p. 1023–1044].

1 Introduction

Fullerenes are cage polyhedral carbon molecules such that all faces are pentagons and hexagons. The icosahedral C_{60} , known as the *buckminsterfullerene*, was predicted by Osawa [14], and discovered by Kroto *et al.* [10]. A fullerene with no adjacent pentagons is said to satisfy the *isolated pentagon rule* (IPR), and is called an *IPR fullerene*. The buckminsterfullerene is the smallest IPR fullerene, and all the stable fullerenes are IPR fullerenes.

Various properties of fullerene molecules can be studied using mathematical tools and results. Thus, *fullerene graphs* were defined as 3-connected cubic plane graphs with all faces of size 5 or 6. Such graphs are suitable models for fullerene molecules: carbon atoms are represented by vertices of the graph, whereas edges represent bonds between adjacent atoms. It is known that there exists a fullerene graph on n vertices for every even

*Research supported by the bilateral Czech-Slovenian project MEB090805.

[†]Institute for Theoretical Computer Science (ITI) and Department of Mathematics, University of West Bohemia, Univerzitní 8, 306 14 Plzeň, Czech Republic. E-mail: kaisert@kma.zcu.cz. Supported by project 1M0545 of the Czech Ministry of Education and project GAČR 201/09/0197 of the Czech Science Foundation.

[‡]CNRS, Laboratoire G-SCOP, 46, avenue Félix Viallet, 38031 Grenoble Cedex 1, France. E-mail: stehlik@kam.mff.cuni.cz.

[§]Department of Mathematics, University of Ljubljana, Jadranska 19, 1111 Ljubljana, Slovenia.

$n \geq 20$, $n \neq 22$. See the monograph of Fowler and Manolopoulos [4] for more information on fullerenes.

Since carbon atoms in stable molecules are tetravalent, one of the three σ -bonds of every atom in a fullerene molecule should be augmented by a π -bond to create a carbon-carbon double bond. A Lewis structure that localises the σ - and π -bonds to at most two atoms per bond is called a *Kekulé structure*. It corresponds to the notion of perfect matchings in fullerene graphs: a *perfect matching* M in a graph G is a set of edges of G such that every vertex of G is incident with precisely one edge of M . It has been shown [1, 7] that fullerene graphs have exponentially many perfect matchings.

Let M be a perfect matching in a fullerene graph G . An even cycle of G is *M -alternating* if exactly half its edges are in M . A set of disjoint M -alternating hexagons of G is called a *resonant pattern*. A fullerene graph is *k -resonant* if any i disjoint hexagons form a resonant pattern, for $0 \leq i \leq k$.

Ye *et al.* [17] proved that every fullerene graph is 1-resonant, and that every *leapfrog* fullerene graph is 2-resonant (see [17] for the definition). The class of leapfrog fullerene graphs satisfies the IPR; this led Ye *et al.* [17, Open Problem 3.5] to ask whether every IPR fullerene graph is 2-resonant. Here we answer this question in the affirmative by proving the following result.

Theorem 1. *Every IPR fullerene graph is 2-resonant.*

The theorem is best possible in the sense that no fullerene graph on more than sixty vertices is 3-resonant, as pointed out by Ye *et al.* [17].

2 Notation and preliminary results

Given a graph G , we let $V(G)$ and $E(G)$ be its vertex set and edge set, respectively. For $H \subseteq V(G) \cup E(G)$, we let $G - H$ be the graph obtained from G by removing the elements in H . Let $X, X' \subset V(G)$. We define $G[X]$ to be the subgraph of G induced by the vertices in X . We let $\partial(X)$ be the set of edges of G with exactly one end-vertex in X ; if $F = G[X]$ we may also write $\partial(F)$ for $\partial(X)$. We denote by $E(X, Y)$ the set of edges with one endvertex in X and the other endvertex in Y ; if $X = Y$ we simply write $E(X)$. We denote $|E(X, Y)|$ and $|E(X)|$ by $e(X, Y)$ and $e(X)$, respectively.

We say that two faces of a plane embedding of G are *adjacent* if they share an edge. Vertices and edges contained in the boundary of a face F are said to *belong to* F or to *be on* F . Pentagonal and hexagonal faces are referred to simply as *pentagons* and *hexagons*.

A graph G is *factor-critical* if $G - v$ has a perfect matching, for every vertex $v \in V(G)$. It is easy to see that every factor-critical graph is connected,

has an odd number of vertices, and has minimum degree two unless it consists of a single vertex. Lovász [12] characterised factor-critical graphs as those having an ‘odd ear decomposition’ (see also [13, 15]). The definition of factor-criticality implies that a factor-critical graph is 2-edge-connected. Using the fact that fullerenes are cubic graphs, we obtain the following slightly stronger conclusion:

Observation 2. *Every non-trivial factor-critical subgraph of a fullerene graph is 2-connected.*

The following theorem may be viewed as a strengthening of Tutte’s 1-factor theorem [16]. It is a simple corollary of the Gallai-Edmonds decomposition theorem [3, 5, 6] (see also [13, 15]).

Theorem 3. *A graph G has no perfect matching if and only if there exists a subset $A \subseteq V(G)$ such that:*

- (i) *every component of $G - A$ has a perfect matching or is factor-critical, and*
- (ii) *at least $|A| + 1$ components of $G - A$ are factor-critical.*

A k -edge-cut of G is a set Y of k edges of G for which there exists a set $X \subset V(G)$ such that $Y = \partial(X)$. A k -edge-cut Y is *cyclic* if at least two components of $G - Y$ contain a cycle. The graph G is *cyclically k -edge-connected* if $|E(G)| > k$, and it contains no cyclic edge-cuts consisting of fewer than k edges.

Došlić [2] proved that fullerene graphs are cyclically 5-edge connected. Kardoš and Škrekovski [9] characterised the cyclic 5- and 6-edge-cuts in fullerene graphs; the cyclic 5-edge-cuts were characterised independently by Kutnar and Marušič [11]. Furthermore, Kardoš *et al.* [8] characterised the cyclic 7-edge-cuts in fullerene graphs. The characterisations together imply the following three lemmas (see Figures 1–3 for illustrations). The proofs are omitted, as they amount to verifying which configurations satisfy the IPR and are factor-critical, resp. have a perfect matching. For a subgraph F of a graph G , the symbol F^c stands for $G - V(F)$.

Lemma 4. *Let F be a factor-critical subgraph of an IPR fullerene graph. Then one of the following assertions holds.*

- (i) *$|\partial(F)| = 3$ and F or F^c consists of a single vertex;*
- (ii) *$|\partial(F)| = 5$ and F or F^c is a pentagon;*
- (iii) *$|\partial(F)| = 7$ and F or F^c consists of a pentagon and an adjacent hexagon;*

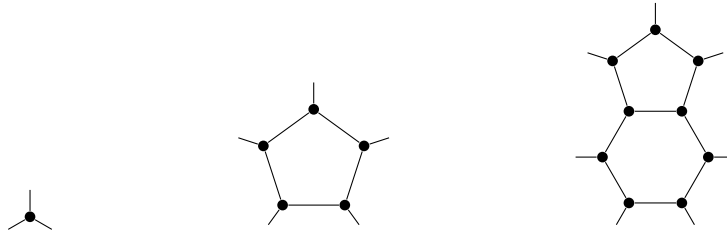


Figure 1: Cases (i)–(iii) in Lemma 4.

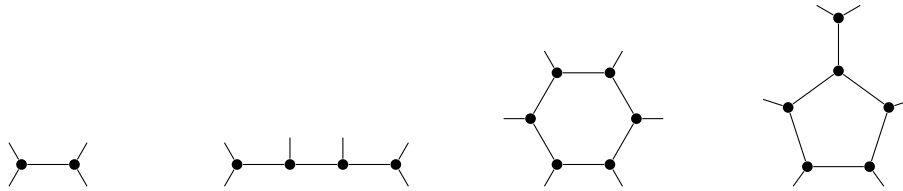


Figure 2: Case (i) and the three subcases of case (ii) in Lemma 5.

(iv) $|\partial(F)| \geq 9$.

Lemma 5. *Let F be a subgraph with a perfect matching of an IPR fullerene graph. Then one of the following assertions holds.*

- (i) $|\partial(F)| = 4$ and F or F^c is isomorphic to K_2 ;
- (ii) $|\partial(F)| = 6$ and F or F^c is a path of length 3, a hexagon, or a pentagon with a pendant edge;
- (iii) $|\partial(F)| \geq 8$.

We will also need to extend the characterisations of Lemmas 4 and 5 to subgraphs F with $|\partial(F)| \leq 6$ that are not necessarily factor-critical and may have no perfect matching. There are only two additional cases (cf. Figure 3):

Lemma 6. *Let F be a subgraph of an IPR fullerene with $|\partial(F)| \leq 6$. Then F or F^c is either a path of length at most 3, a $K_{1,3}$, a cycle of length 5 or 6, or a pentagon with a pendant edge.*



Figure 3: The two new cases in Lemma 6. The remaining possibilities are the first two graphs in Figure 1 and the graphs in Figure 2.

3 The proof of Theorem 1

Before embarking on the proof of Theorem 1, let us give a brief outline of the main argument. Using Theorem 3, we show that deleting any pair of disjoint hexagons from a fullerene graph results in a graph which is ‘almost bipartite’, in the sense that deleting a prescribed set of vertices and edges gives a bipartite graph. Then, using Lemmas 4 and 5 we show that, apart from one special case, there is always at least one pentagon which does not intersect this prescribed set; a contradiction. The remainder of the proof is then dedicated to ruling out the special case.

Proof of Theorem 1. Let G be an IPR fullerene graph, let H_1 and H_2 be two disjoint hexagons of G , and set $H := V(H_1 \cup H_2)$. Suppose $G - H$ does not have a perfect matching. Then Theorem 3 guarantees the existence of a set A of vertices of $G - H$ such that $G - (H \cup A)$ contains more than $|A|$ factor-critical components, and all other components of $G - (H \cup A)$ have a perfect matching. We denote the factor-critical components of $G - (H \cup A)$ by \mathcal{D} , and the components of $G - (H \cup A)$ with a perfect matching by \mathcal{C} . The subset of *non-trivial* factor-critical components of $G - (H \cup A)$ is denoted by \mathcal{D}^* . The union of the vertex sets of the components of \mathcal{C} , \mathcal{D} and \mathcal{D}^* is denoted by C , D and D^* , respectively, and we set $D_0 := D - D^*$ (see Figure 4).

The number of vertices of G being even, $|\mathcal{D}|$ and $|A|$ have the same parity. Since $|\mathcal{D}| > |A|$, we deduce that

$$|\mathcal{D}| \geq |A| + 2. \tag{1}$$

The number of edges leaving $A \cup H$ may be expressed as

$$|\partial(A \cup H)| = |\partial(H)| + 3|A| - 2e(A, H) - 2e(A). \tag{2}$$

Now set

$$s(\mathcal{D}) := \sum_{F \in \mathcal{D}} (|\partial(F)| - 3)/2$$

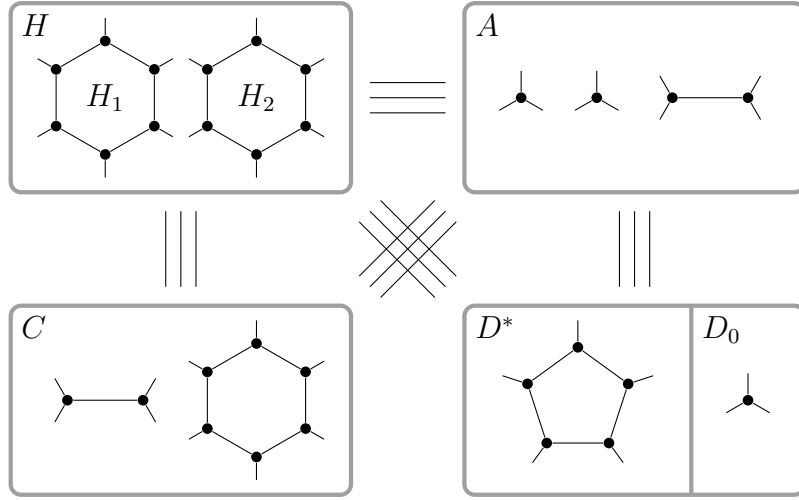


Figure 4: The partition of $V(G)$ into H , A , C , D^* and D_0 .

and

$$s(\mathcal{C}) := \sum_{F \in \mathcal{C}} (|\partial(F)| - 4)/2.$$

By Lemmas 4 and 5, the number of edges leaving $C \cup D$ may be expressed as

$$|\partial(C \cup D)| = 3|\mathcal{D}| + 4|\mathcal{C}| + 2(s(\mathcal{C}) + s(\mathcal{D})). \quad (3)$$

Since $\partial(A \cup H) = \partial(C \cup D)$, (1), (2) and (3) imply that

$$2|\mathcal{C}| + e(A, H) + e(A) + s(\mathcal{C}) + s(\mathcal{D}) \leq \frac{1}{2}|\partial(H)| - 3 \leq 3. \quad (4)$$

Let \mathcal{P} be the set of pentagons of G ; it follows from Euler's formula that $|\mathcal{P}| = 12$. If $X \subseteq V(G)$ and $Y \subseteq E(G)$, let $p(X) := |\{P \in \mathcal{P} \mid V(P) \cap X \neq \emptyset\}|$ and $p(Y) := |\{P \in \mathcal{P} \mid E(P) \cap Y \neq \emptyset\}|$.

Observe that every edge of the graph $G - (H \cup C \cup D^* \cup E(A))$ has one endvertex in A and the other in D_0 ; in particular, $G - (H \cup C \cup D^* \cup E(A))$ is bipartite. This means that all twelve pentagons of G must contain a vertex in $H \cup C \cup D^*$ or an edge in $E(A)$. By the IPR, $p(H) \leq 6$, so at least six pentagons disjoint from H must have a vertex in $C \cup D^*$ or an edge in $E(A)$. In particular,

$$p(E(A)) + p(C) + p(D^*) \geq 6. \quad (5)$$

Claim 1. *There exists a component $F^* \in \mathcal{D}^*$ such that $|\partial(F^*)| = 9$.*

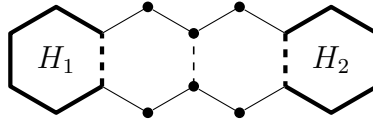


Figure 5: A link of size 3. The hexagons H_1 and H_2 are drawn bold, the edges of the link are dashed.

By (4), $|\mathcal{C}| \leq 1$. If $|\mathcal{C}| = 1$, then $e(A) + s(\mathcal{D}) + s(\mathcal{C}) \leq 1$ by (4). By the IPR, $p(E(A)) \leq e(A) \leq 1$. As $s(\mathcal{D}) \leq 1$, Lemma 4 and the IPR imply that $p(D^*) \leq 1$. Finally, by Lemma 5 and the IPR, $p(C) \leq s(C) + 2 \leq 3$. Hence $p(E(A)) + p(C) + p(D^*) \leq e(A) + s(\mathcal{D}) + s(\mathcal{C}) + 2 < 6$, contradicting (5).

Hence we may assume that $\mathcal{C} = \emptyset$. In particular, $p(C) = s(C) = 0$, and by (4) $e(A) + s(\mathcal{D}) \leq 3$. This implies that $|\partial(F)| \leq 9$ for every $F \in \mathcal{D}^*$. If $|\partial(F)| < 9$ for every $F \in \mathcal{D}^*$, then by Lemma 4 and the IPR, $p(D^*) \leq s(\mathcal{D}) + 1$. Hence $p(E(A)) + p(C) + p(D^*) \leq e(A) + s(\mathcal{D}) + 1 < 6$, contradicting (5). Therefore $|\partial(F^*)| = 9$ for some component $F^* \in \mathcal{D}^*$, as claimed. \triangle

By (4), F^* is the unique non-trivial component of \mathcal{D} , $\mathcal{C} = \emptyset$, $e(A) = 0$ (so A is an independent set), $e(A, H) = 0$ and $|\partial(H)| = 12$ (in particular, there are no edges between the hexagons H_1 and H_2). As F^* is non-trivial and factor-critical, it is 2-connected by Observation 2. In particular, it has minimum degree 2.

Let us now fix a planar embedding of G . If e, f are edges of G , we will say that e is *opposite* f if they both belong to the boundary of some hexagonal face and no edge of this boundary is incident with both e and f . A *link* is a sequence $L = (e_1, \dots, e_k)$ of edges, where

- (1) e_{i+1} is opposite e_i for $i = 1, \dots, k - 1$,
- (2) e_1 belongs to H_1 and e_k belongs to H_2 .

(See Figure 5 for an illustration.) The *size* $|L|$ of such a link is defined as k . The following observation is immediate:

Observation 7. *Any two links in G are disjoint.*

We continue with a simple structural lemma.

Lemma 8. *G contains no two links L_1, L_2 with $|L_1| + |L_2| \leq 6$.*

Proof. Suppose such L_1 and L_2 exist. We claim that their union (as sets) forms an edge-cut of G . This is easy to see from the duality between edge-cuts in G and cycles in the dual graph. For a direct argument, note that one

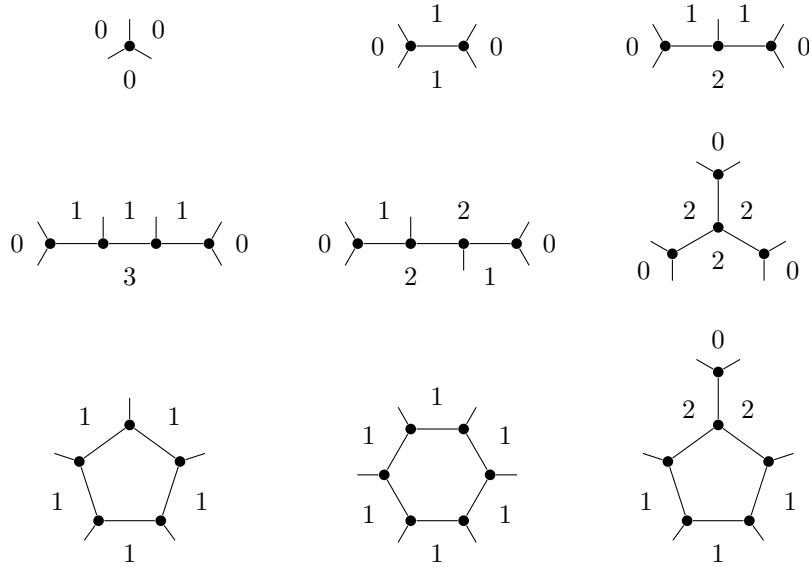


Figure 6: The values of $\ell(F)$ in the proof of Lemma 8 (shown as the numbers inside the faces adjacent to G_1).

can draw a closed curve C in the plane which only intersects the edges in $L_1 \cup L_2$ (connecting each pair of opposite edges by a line segment through the face containing them, and proceeding similarly on H_1 and H_2). In addition, the subgraphs G_1 and G_2 into which G is divided by C can easily be seen to be connected, which shows that $L_1 \cup L_2$ is an edge-cut.

Lemma 6 gives us the possibilities for G_1 or G_2 . Let us assume that G_1 is one of the subgraphs listed in the lemma. Let N be the set of faces of G which intersect both G_1 and G_2 . For each face F in N , let $\ell(F)$ be the number of edges of G_1 that belong to F . By the definition of a link, we have $\ell(F) = 2$ for each $F \in N \setminus (H_1 \cup H_2)$. However, considering the possibilities for G_1 of Lemma 6 (with two different embeddings in the case of the path of length 3) we find that in each case, at least three values of $\ell(F)$ are different from 2. See Figure 6 for an illustration. This contradiction concludes the proof. \square

The following lemma is the core of our argument.

Lemma 9. *Let $j \in \{1, 2\}$. Assume that a, b are adjacent vertices of H_j and M is the face of G whose boundary includes the edge ab and which is different from H_j . Let a', b' be the neighbours of a and b (respectively) not in H_j . If neither aa' nor bb' is an edge of $\partial(F^*)$, then M is either a pentagonal face,*

or a hexagonal face intersecting both H_1 and H_2 .

Proof. Suppose that M has length 6, so its boundary is $a'abb'xy$, where $x, y \in V(G)$. Since a' and b' belong neither to $A \cup V(H_1 \cup H_2)$ nor to F^* (by assumption), we have $a', b' \in D_0$. Thus, x and y are contained in $A \cup V(H_1 \cup H_2)$. Since $E(A \cup V(H_1) \cup V(H_2)) = E(V(H_1)) \cup E(V(H_2))$, both x and y must be contained in the same hexagon H_k (for some $k = 1, 2$). However, if $x, y \in V(H_j)$, then it is easy to find a cycle of length at most 4 in G , contradicting the fact that G is cyclically 5-edge-connected. Thus, M intersects both H_1 and H_2 which proves the claim. \square

We want to be able to bound the number of edges of $\partial(F^*)$ ‘near’ H_1 and H_2 . We will make use of the following lemma.

Lemma 10. *Let F be a face of G incident with H_j ($j = 1, 2$). Then the intersection of $\partial(F^*)$ with the boundary of F is either empty, or it is a matching of size 2 which does not form a pair of opposite edges. If, in addition, F is a pentagon, then both edges of the matching have one endvertex on H_j .*

Proof. If the intersection of $\partial(F^*)$ with the boundary of F is non-empty, then it has to be a matching $\{vw, xy\}$, by Observation 2 and the fact that the intersection of a cycle and an edge cut is even. If both vw and xy have a vertex in $V(H_j)$, then vw, xy are not opposite, as required.

So suppose now that one of the edges vw, xy has no vertex in $V(H_j)$; without loss of generality say it is vw . If v is adjacent to a vertex of H_j , then $v \in D_0$ and $w \in D^*$, contradicting the fact that there are no edges between D_0 and D^* . Similarly for w . Therefore v, w are not adjacent to any vertex of $V(H_j)$, so M is a hexagon and vw, xy are not opposite. \square

A *cluster* at H_j ($j = 1, 2$) is a sequence $Q = (e_1, \dots, e_k)$ such that:

- each e_i is in $\partial(F^*)$,
- e_i and e_{i+1} are on the same face ($1 \leq i \leq k - 1$),
- for $2 \leq i \leq k - 1$, each e_i has exactly one endvertex on H_j , while e_1 and e_k are not incident with H_j .

Examples of clusters are shown in Figure 7. The size $|Q|$ of Q is the number of its edges. Observe that $|Q| \geq 3$.

Lemma 11.

- (i) *Clusters at H_j are pairwise disjoint (for $j = 1, 2$).*

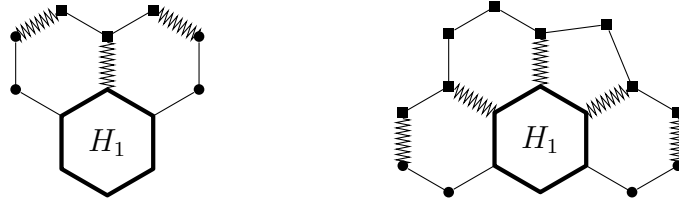


Figure 7: Examples of clusters at H_1 . Left: a cluster of size 3. Right: a cluster of size 5. The edges of the clusters are shown as zigzag lines, those of H_1 are drawn bold. Squares represent the vertices of F^* .

- (ii) G contains at most one edge e that is contained in a cluster at H_1 and a cluster at H_2 . If e exists, then it is contained in a link of size 3.

Proof. (i) If an edge e belongs to two distinct clusters C, C' at H_j , then it is clearly not incident with H_j . We may thus assume that C and C' are of the form $C = (e, e_2, \dots, e_k)$ and $C' = (e, e'_2, \dots, e'_l)$. By the definition of a cluster and Lemma 10, e belongs to hexagonal faces M and M' whose boundaries contain e_2 and e'_2 , respectively. However, e belongs to the boundary of only one face adjacent to H_j , so $M = M'$ and this face boundary contains e, e_2 and e'_2 , a contradiction to Lemma 10.

(ii) Assume an edge e is contained in a cluster C_1 at H_1 and C_2 at H_2 . By Lemma 10 and the fact that there are no edges between H_1 and H_2 , it is easy to check that e is incident with neither H_1 nor H_2 . By Lemma 10 e is opposite an edge of H_1 and opposite an edge of H_2 , so it is contained in a link of size 3. The uniqueness of e follows by Lemma 8. \square

We shall now use Lemma 9 to lower-bound the number of edges of $\partial(F^*)$ incident with each H_j . Let us call an edge e of G *free* if $e \notin \partial(F^*)$.

Lemma 12. *Let $j \in \{1, 2\}$.*

- (i) $\partial(H_j)$ does not contain five free edges.
- (ii) If $\partial(H_j)$ contains three consecutive¹ free edges, then there is a link of size 2 in G .
- (iii) There is exactly one cluster at H_j , and its size is 4 or 5.

Proof. (i) Any set of five free edges in $\partial(H_j)$ contains four consecutive pairs of free edges. By Lemma 9, each of these pairs belongs to the boundary

¹‘Consecutive’ refers to the natural clockwise ordering of $\partial(H_j)$.

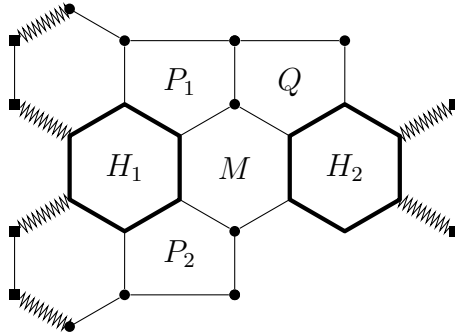


Figure 8: An illustration to the final argument in the proof of Theorem 1 (the case when Q is adjacent to P_1). The conventions for the vertices and edges are the same as in Figure 7.

of either a pentagonal face, or a hexagonal face incident with H_1 and H_2 . Lemma 8 implies that only one of the four pairs is of the latter type. It follows that $\partial(H_j)$ contains three edges e, e', e'' such that each of the pairs e, e' and e', e'' belongs to a pentagon. Since this would violate the IPR, $\partial(H_j)$ does not contain five free edges.

(ii) This is a direct consequence of Lemma 9 and the IPR.

(iii) Suppose that $\partial(H_j) \subseteq \partial(F^*)$, or that H_j is incident with two or more clusters. In both cases there are at most three edges of $\partial(F^*)$ left for the clusters at H_{3-j} . Hence there is only one cluster (of size 3) at H_{3-j} and five edges of $\partial(H_{3-j})$ are free, contradicting (i). The claim follows. \square

We now use Lemma 12(iii) to conclude the analysis. Assume first that there is a cluster of size 5 at both H_1 and H_2 . These clusters have nonempty intersection as $|\partial(F^*)|$ is only 9. By Lemma 11(ii), G contains a link of size 3. On the other hand, observe that the set $\partial(H_1)$ contains three consecutive free edges, which by Lemma 12(ii) is only possible if there is a link of size 2 in G . The existence of these two links contradicts Lemma 8.

We may thus assume that the cluster at H_1 has size 4. Since there are four consecutive free edges $e_1, \dots, e_4 \in \partial(H_1)$, Lemma 12(ii) implies the existence of a link L of size 2 in G . Figure 8 depicts the situation. Let M be the (hexagonal) face whose boundary contains both edges of L . Lemma 8 and the fact that there are no edges between the hexagons H_1 and H_2 implies that M is the unique face intersecting both H_1 and H_2 . Therefore, by Lemma 9 and the IPR, e_2 and e_3 belong to M , while the pairs e_1, e_2 and e_3, e_4 belong to pentagonal faces (P_1 and P_2 , respectively) adjacent to H_1 and M .

The cluster at H_2 has size 4 or 5, which means that $\partial(H_2)$ contains at least three consecutive free edges. By another application of Lemma 9 and

the IPR, two of them must belong to a pentagonal face Q adjacent to M and H_2 . However, it is easy to see that Q is adjacent to either P_1 or P_2 . This contradiction to the IPR completes the proof of Theorem 1. \square

References

- [1] M. Chudnovsky and P. Seymour. Perfect matchings in planar cubic graphs. *Combinatorica*. To appear.
- [2] T. Došlić. Cyclical edge-connectivity of fullerene graphs and $(k, 6)$ -cages. *J. Math. Chem.*, 33(2):103–112, 2003.
- [3] J. Edmonds. Paths, trees, and flowers. *Canad. J. Math.*, 17:449–467, 1965.
- [4] P. W. Fowler and D. E. Manolopoulos. *An Atlas of Fullerenes*. Oxford University Press, Oxford, 1995.
- [5] T. Gallai. Kritische Graphen. II. *Magyar Tud. Akad. Mat. Kutató Int. Közl.*, 8:373–395, 1963.
- [6] T. Gallai. Maximale Systeme unabhängiger Kanten. *Magyar Tud. Akad. Mat. Kutató Int. Közl.*, 9:401–413, 1964.
- [7] F. Kardoš, D. Král', D. Miškuf, and J.-S. Sereni. Fullerene graphs have exponentially many perfect matchings. *J. Math. Chem.*, 46(2):443–477, 2009.
- [8] F. Kardoš, M. Krnc, B. Lužar, and R. Škrekovski. Cyclic 7-edge-cuts in fullerene graphs. *J. Math. Chem.*, 47(2):771–789, 2010.
- [9] F. Kardoš and R. Škrekovski. Cyclic edge-cuts in fullerene graphs. *J. Math. Chem.*, 44(1):121–132, 2008.
- [10] H. W. Kroto, J. R. Heath, S. C. O'Brien, R. F. Curl, and R. E. Smalley. C_{60} : Buckminsterfullerene. *Nature*, 318:162–163, 1985.
- [11] K. Kutnar and D. Marušič. On cyclic edge-connectivity of fullerenes. *Discrete Applied Math.*, 156(10):1661–1669, 2008.
- [12] L. Lovász. A note on factor-critical graphs. *Studia Sci. Math. Hungar.*, 7:279–280, 1972.

- [13] L. Lovász and M. D. Plummer. *Matching theory*, volume 121 of *North-Holland Mathematics Studies*. North-Holland Publishing Co., Amsterdam, 1986. *Annals of Discrete Mathematics*, 29.
- [14] E. Osawa. Superaromaticity. *Kagaku (Kyoto)*, 25:854–863, 1970.
- [15] A. Schrijver. *Combinatorial optimization: Polyhedra and efficiency*, volume 24 of *Algorithms and Combinatorics*. Springer-Verlag, Berlin, 2003.
- [16] W. T. Tutte. The factorization of linear graphs. *J. London Math. Soc.*, 22:107–111, 1947.
- [17] D. Ye, Z. Qi, and H. Zhang. On k -resonant fullerene graphs. *SIAM J. Discrete Math.*, 23(2):1023–1044, 2009.